

# Grammar Implementation with TAG

## TuLiPA and the lexicon

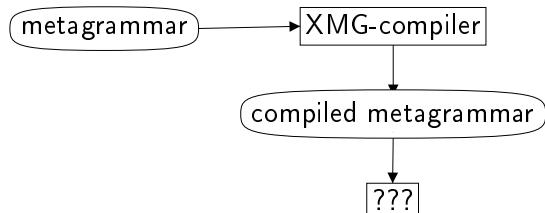
Timm Lichte

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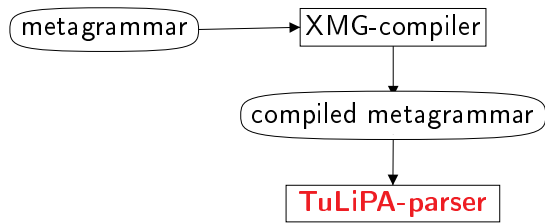
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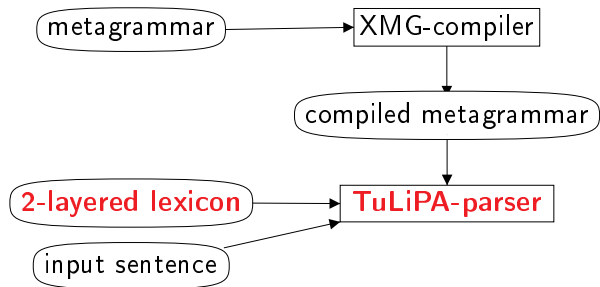
# The situation



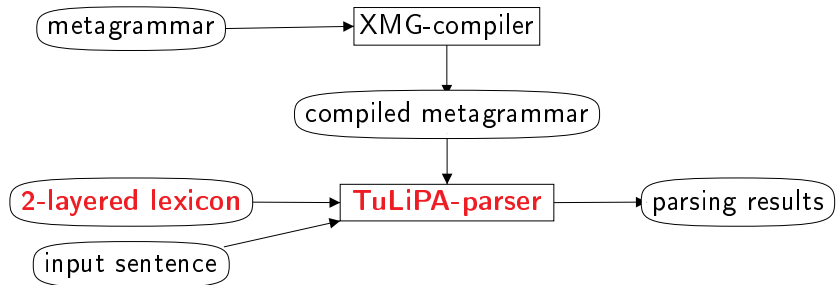
# The missing pieces



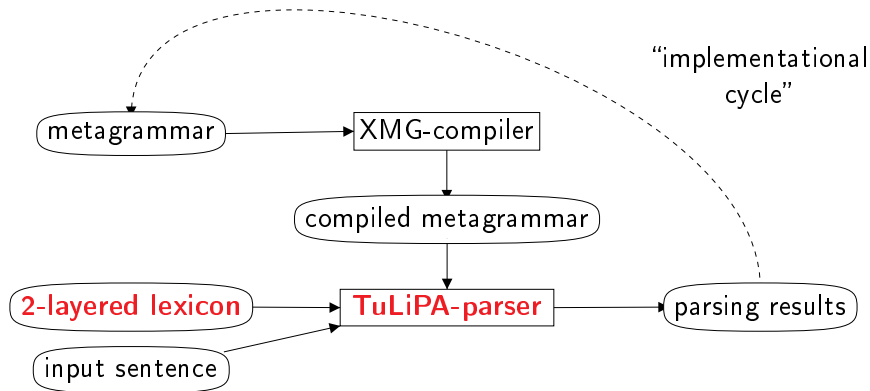
# The missing pieces



# The missing pieces



# The missing pieces



TuLiPA uses Range Concatenation Grammar (RCG) as a pivot formalism.

## Components:

- 1 TAG-to-RCG converter (on-line)
- 2 RCG parser → RCG derivation forest → TAG derivation forest
- 3 Parse viewer (derived tree, derivation tree, dependency view, semantic representation)

## Availability of TuLiPA:

written in Java and released under the GNU GPL  
(<http://sourcesup.cru.fr/tulipa/>)

## Range Concatenation Grammar (RCG)

$$A(X_1 X_2, X_3) \rightarrow B(X_2) C(X_1, X_3)$$

$A, B, C$ : predicates

$X_1, X_2, X_3$ : range variables (instantiated with substrings of the string)

### Example:

$$S(X_1 X_2) \rightarrow NP(X_1) VP(X_2)$$

$$VP(X_3 X_4) \rightarrow V(X_3) NP(X_4)$$

$$NP(\text{Peter}) \rightarrow \epsilon$$

$$NP(\text{Susan}) \rightarrow \epsilon$$

$$V(\text{loves}) \rightarrow \epsilon$$



## Simple RCG

- **non-combinatorial:** each argument on the RHS of a clause consists of a unique variable!

NOT:  $A(XY, Z) \rightarrow B(Y) C(XZ)$

- **linear:** each variable appears at most once in the LHS or RHS of a clause!

NOT:  $A(XYX, Z) \rightarrow B(Y) C(X, Z)$

- **non-erasing:** every variable in the LHS of a clause is also in its RHS!

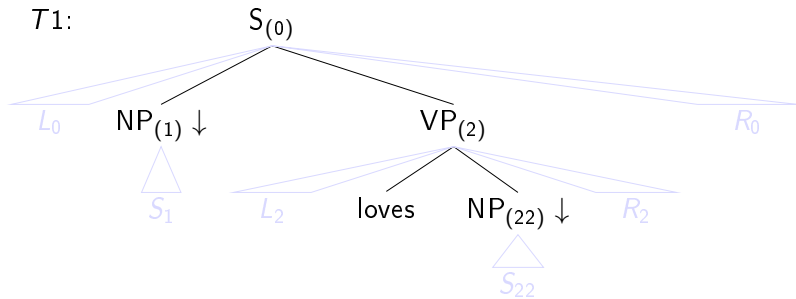
NOT:  $A(XY, Z) \rightarrow B(Y) C(X)$

RCG is useful, because:

- it has attractive formal properties (polynomially parsable, full expressive power of MCS-languages);
- there exist parsing algorithms.

⇒ Parser can be reused for other mildly context-sensitive formalisms!

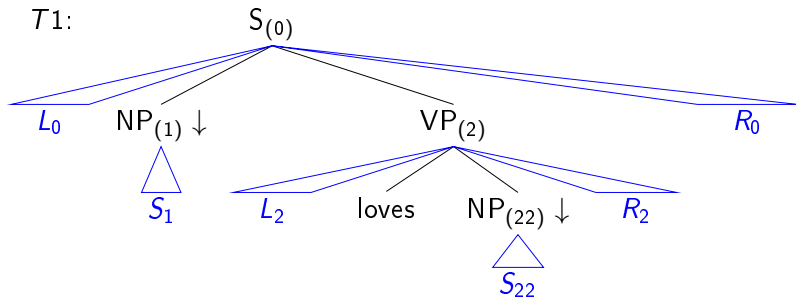
NB: RCG properly includes MCS. We use a restricted RCG, called *simple RCG*, that is included in MCS.



$T_1(L_0 S_1 L_2 \text{ loves } S_{22} R_2 R_0) \rightarrow$   
 $adj_{S_0}(L_0, R_0) \text{ sub}_{NP_1}(S_1) \text{ adj}_{VP_2}(L_2, R_2) \text{ sub}_{NP_{22}}(S_{22})$

$adj_{S_0}(X, Y) \rightarrow T_2(X, Y)$

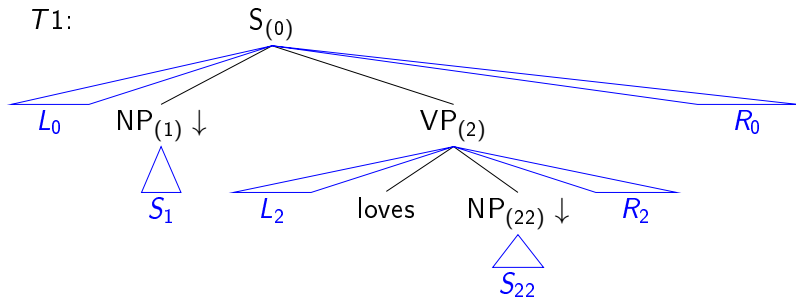
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$T_1(L_0 S_1 L_2 \text{ loves } S_{22} R_2 R_0) \rightarrow$   
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...



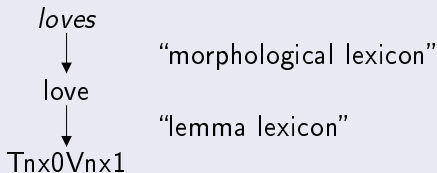
$$T_1(L_0 S_1 L_2 \text{ loves } S_{22} R_2 R_0) \rightarrow$$

$$adj_{S_0}(L_0, R_0) \text{ sub}_{NP_1}(S_1) \text{ adj}_{VP_2}(L_2, R_2) \text{ sub}_{NP_{22}}(S_{22})$$

$$adj_{S_0}(X, Y) \rightarrow T_2(X, Y)$$

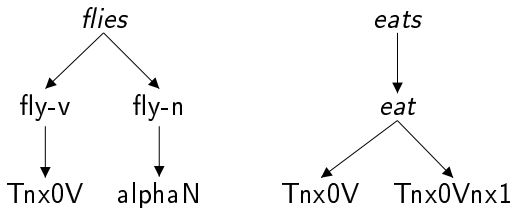
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## A 2-layered lexicon



The morphological lexicon and the lemma lexicon are compiled using `lexConverter` (as a part of `LEX2ALL`).

## A 2-layered lexicon - 2 sources for ambiguity



# A 2-layered lexicon

## Morphological lexicon

maps an (inflected) token to some base form (= lemma), while preserving morphological information in a feature structure.

loves	love	[pos=v; num=sing; pers=3;]
Peter	Peter	[pos=n; num=sing; pers=3; case=nom acc;]

## Interface with tree templates:

Feature unification during lexical insertion



## Lemma lexicon

maps a lemma onto tree tuple families, while also containing selectional restrictions (e.g., case assignment).

```
*ENTRY: love
*CAT: v
*SEM:
*ACC: 1
*FAM: Tnx0Vnx1
*FILTERS: []
*EX:
*EQUATIONS:
NParg1 -> case = nom
NParg2 -> case = acc
*COANCHORS:
```

### Interface with tree templates:

EQUATIONS → nodes of tree templates (via name property)

FILTERS → tree templates (via interface expressions)