

## Example: Shift-Reduce Parsing

Philipp Stumpe

Consider the CFG  $G = \langle N, T, P, S \rangle$  with  $N = \{S, A, B, C, D\}$ ,  $T = \{a, b, c, d\}$  and

$$P = \{ \begin{array}{l} S \rightarrow aA \\ A \rightarrow BCD \\ B \rightarrow BS|b \\ C \rightarrow c \\ D \rightarrow d \end{array} \},$$

and the input  $w = ababcdcd$ .

1. Give the trace of the shift-reduce parsing without deduction rules for this input.

input position	id	stack	remaining input	operation
0	1	$\epsilon$	ababcdcd	-
1	2	a	babcdcd	shift (1)
2	3	ab	abcdcd	shift (2)
2	4	aB	abcdcd	reduce (3), $b \rightarrow B$
3	5	aBa	bcdcd	shift (4)
4	6	aBab	cdcd	Shift (5)
4	7	aBaB	cdcd	Reduce (6), $b \rightarrow B$
5	8	aBaBc	dcd	Shift (7)
5	9	aBaBC	dcd	Reduce (8), $c \rightarrow C$
6	10	aBaBCd	cd	Shift (9)
6	11	aBaBCD	cd	Reduce (10), $d \rightarrow D$
6	12	aBaA	cd	Reduce (11), $BCD \rightarrow A$
6	13	aBS	cd	Reduce (12), $aA \rightarrow S$
6	14	aB	cd	Reduce (13), $BS \rightarrow B$
7	15	aBc	d	Shift (14)
7	16	aBC	d	Reduce (15), $c \rightarrow C$
8	17	aBCd	$\epsilon$	Shift (16)
8	18	aBCD	$\epsilon$	Reduce (17), $d \rightarrow D$
8	19	aA	$\epsilon$	Reduce (18), $BCD \rightarrow A$
8	20	S	$\epsilon$	Reduce (19), $aA \rightarrow S$

2. If  $w \in L(G)$ , read off the rightmost derivation for  $w$ .

$S \rightarrow aA \rightarrow aBCD \rightarrow aBCd \rightarrow aBcd \rightarrow aBScd \rightarrow aBaAcd \rightarrow aBaBCDcd \rightarrow aBaBCdcd \rightarrow aBaBcdcd \rightarrow aBabcdcd \rightarrow ababcdcd$