# Mildly Context-Sensitive Grammar Formalisms:

## Multicomponent TAG

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## Motivation (1)

- Multicomponent Tree Adjoining Grammars (MCTAGs) were first introduced in [Joshi et al., 1975] as *simultaneous TAGs*, later redefined as *multicomponent TAGs (MCTAGs)* in [Weir, 1988, Joshi, 1985].
- The underlying linguistic motivation is the idea to separate the contribution of a lexical item into several components.
- In each derivation step, a new set is picked and all trees from the set are added simultaneously, i.e., they are attached (by substitution or adjunction) to different nodes in the already derived tree.

Grammar Formalisms	1	MCTAG	Grammar Formalisms	3	MCTAG
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			Motivation (2)		
			(1) which painting <sub>i</sub> did you s	see a picture of	$t_i$



- 1. Motivation
- 2. Different types of MCTAG
- 3. Semilinearity of MCTALs

NP S aux S

did NP



Det

NP

which painting

Grammar Formalisms

MCTAG

## Motivation (3)

Constructions that require multicomponents:

- Extraction out of complex NPs [Kroch, 1989], stranding phenomena, in particular "picture-NPs":
  (2) which castle did you paint a picture of?
- Subject-aux inversion in raising questions [Frank, 2008]
- (3) Does John seem to annoy you?
- Scrambling in German [Rambow, 1994]
- (4) dass den Kühlschrank niemand zu reparieren versprochen hat

Without loss of generality, we can assume that  $\mathcal{A}$  is a partition of

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## Different types of MCTAG (2)

**Definition 2 (MCTAG derivation)**  $\gamma \Rightarrow \gamma'$  is a derivation step in G iff there is an instance  $\{\gamma_1, \ldots, \gamma_n\}$  of an elementary tree set in  $\mathcal{A}$  and there are pairwise different nodes  $v_1, \ldots, v_n$  in  $\gamma$  such that  $\gamma' = \gamma[v_1, \gamma_1] \ldots [v_n, \gamma_n].$ 

As in TAG, a derivation starts from an initial tree and in the end, in the final derived tree, all leaves must have terminal labels (or the empty word) and there must not be any OA constraints left.

Grammar Formalisms	5	MCTAG	Grammar Formalisms	7	MCTAG
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Different types of M	CTAG (1)		Different types of MCTAG (3)		
Definition 1 (MCTA	G) An MCTAG is a tup	ble	An MCTAG is called		
$G = \langle N, T, S, I, A, f_{OA}, f_{SA}, \mathcal{A} \rangle$ such that:			• <i>tree-local</i> iff in each derivation step, the nodes the new trees attach to belong to the same elementary tree.		
• $G_{TAG} := \langle N, T, S, I, A, f_{OA}, f_{SA} \rangle$ is a TAG with adjunction constraints, and					
			• <i>set-local</i> iff in each derivation step, the nodes the new trees		
_ ( )	set of subsets of $I \cup A$ , t	he set of elementary	attach to belong to the same elementary tree set.		tree set.
tree sets. <sup>a</sup>			• non-local otherwise		

• non-local otherwise.

Usually, the term "MCTAG" without specification of the locality means "set-local MCTAG".

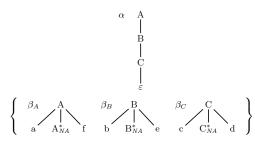
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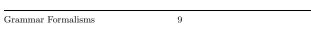
 ${}^{a}P(X)$  is the set of subsets of some set X.

 $I \cup A$ .

## Different types of MCTAG (4)

Set-local MCTAG for  $L_6 = \{a^n b^n c^n d^n e^n f^n \mid n \ge 0\}$ :



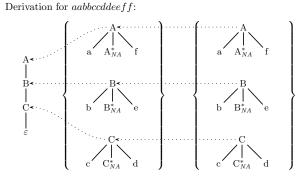


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MCTAG

## Different types of MCTAG (5)



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#### Different types of MCTAG (6)

Tree-local MCTAG and TAG are equivalent since we can precompile the possible adjunctions and substitutions in an elementary tree:

Proposition 1 Tree-local MCTAG are strongly equivalent to TAG.

For a given tree-local MCTAG, a strongly equivalent TAG can be easily constructed adopting corresponding adjunction constraints that enforce the simultaneous adjunctions of all elementary trees from a tree set.

But: the number of elementary trees in the grammar can increase in an exponential way in this construction ( $\Rightarrow$  rather a bad strategy for tree-local MCTAG parsing).

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#### Semilinearity of MCTALs (1)

For tree-local and set-local MCTAG, *derivation trees* can be defined as follows [Weir, 1988]:

- each node is an ordered elementary tree set (the initial tree the derivation starts with is considered as a unary set),
- each edge represents the simultaneous adjunctions of the trees from the daughter tree set to nodes in the trees in the mother tree set; an edge is equipped with a tuple of n node positions where n is the number of trees in the daughter set. Each node position is of the form  $\langle \gamma, p \rangle$  where  $\gamma$  is one of the trees in the mother set and p a position in  $\gamma$ .

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## Semilinearity of MCTALs (2)

Derivation tree of the previous sample derivation

$$\begin{array}{c} \langle \alpha \rangle \\ & & \\ \langle \langle \alpha, \epsilon \rangle, \langle \alpha, 1 \rangle, \langle \alpha, 11 \rangle \rangle \\ \langle \beta_A, \beta_B, \beta_C \rangle \\ & & \\ \langle \langle \beta_A, \epsilon \rangle, \langle \beta_B, \epsilon \rangle, \langle \beta_C, \epsilon \rangle \rangle \\ \langle \beta_A, \beta_B, \beta_C \rangle \end{array}$$

## Semilinearity of MCTALs (4)

Letter-equivalent CFG for our sample MCTAG:

 $S \to \langle \alpha \rangle$  $\langle \alpha \rangle \to \varepsilon$  $\langle \alpha \rangle \rightarrow \langle \beta_A, \beta_B, \beta_C \rangle abcdef$  $\langle \beta_A, \beta_B, \beta_C \rangle \to \varepsilon \quad \langle \beta_A, \beta_B, \beta_C \rangle \to \langle \beta_A, \beta_B, \beta_C \rangle abcdef$ 

We will see later that MCTALs are also in PTIME and consequently mildly context-sensitive.

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Semilinearity of MCTALs (3)			References		
Construction of a letter-equivalent CFG for a given set-local MCTAG: The non-terminals are a start symbol $S$ and the tuples indicating the elementary tree sets.			[Frank, 2008] Frank, R. (2008). Syntax and Itag. Slides of a tutorial at TAG+9, Tübingen.		
• For every unary elementary tree set tuple $\Gamma$ containing an initial tree with root label $S$ where $\Gamma$ contains terminals $a_1 \dots a_m$ , we add $S \to \Gamma a_1 \dots a_m$			[Joshi, 1985] Joshi, A. K. (1985). Tree adjoining grammars: How much contextsensitivity is required to provide reasonable structural descriptions? In Dowty, D., Karttunen, L., and Zwicky, A., editors, <i>Natural Language Parsing</i> , pages 206–250. Cambridge University Press.		
• Whenever a tuple $\Gamma$ allows us to attach the tuples $\Gamma_1, \ldots \Gamma_k$ and the new trees add the terminals $a_1 \ldots a_m$ , we add			[Joshi et al., 1975] Joshi, A. K., Levy, L. S., and Takahashi, M.		

$$\Gamma \to \Gamma_1, \ldots \Gamma_k a_1 \ldots a_m$$

• For every tuple  $\Gamma$  containing neither OA-nodes nor substitution nodes, we add

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$$\Gamma \rightarrow \varepsilon$$

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- System Science, 10:136–163. [Kroch, 1989] Kroch, A. (1989). Asymmetries in long-distance

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