

# Formal Languages and Automata Theory

## Homework 4 (FSA), Abgabe 14.11.2017

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**Exercise 1** (3 points) Convert the NFA defined by

- $\delta(q_0, a) = \{q_0, q_1\}$
- $\delta(q_1, b) = \{q_1, q_2\}$
- $\delta(q_2, a) = \{q_2\}$
- $\delta(q_1, \epsilon) = \{q_1, q_2\}$

with initial state  $q_0$  and final state  $q_2$  into an equivalent DFA.

**Exercise 2** (3 points) In converting NFA to DFA, the number of states may increase substantially. Give upper and lower bounds on the increase in number of states for an  $n$ -state NFA.

**Exercise 3** (4 points) Let  $L$  be any language. Define  $\mathit{even}(w)$  as the string obtained by extracting from  $w$  the letters in even-numbered positions; that is, if  $w = a_1a_2a_3a_4a_5a_6\dots$ , then  $\mathit{even}(w) = a_2a_4a_6\dots$

Corresponding to this, we can define a language  $\mathit{even}(L) = \{\mathit{even}(w) : w \in L\}$ .

Prove that if  $L$  is regular, so is  $\mathit{even}(L)$ .

**Exercise 4** (4 points) Show that if  $L$  is regular, so is  $L^R$ .

**Exercise 5** (4 points) Is it true that for NFA  $M = (Q, \Sigma, \delta, q_0, F)$ , the complement of  $L(M)$  is equal to the set  $\{w \in \Sigma^* : \delta^*(q_0, w) \cap F = \emptyset\}$ ? If so, prove it. If not, give a counterexample.